

*Dissertation Proposal*

# Sensor Deployment Algorithms for Target Positioning Services

Pei-Ling Chiu

Advisor: Prof. Frank Yeong-Sung Lin  
Department of Information Management  
National Taiwan University

## Outline

- Overview
- Research Background
- Research Problems
  - Sensor placement problem for complete coverage and complete discrimination
  - Random deployment problem using adjustable sensing radius
  - Sensor placement problem for differentiated quality of positioning services
  - Energy-efficient sensor networks design
- Conclusion

## Overview

- Motivation
- Research Scope
- Research Problems
- Research Methodologies

## Motivation

- **Wide range of applications in WSNs**
  - Surveillance, target positioning, environment monitoring, health caring, and animal tracking
- **Trend of user-centric services**
  - Location-based services, personalized services
  - Localization capability is one of the most important techniques for supporting these applications

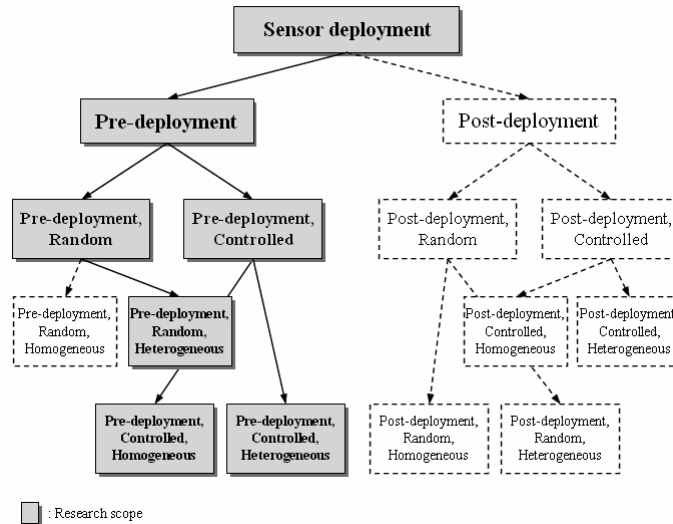
## Motivation

- **Localization Systems**
  - GPS is not useful in indoor
  - Three techniques of indoor location systems: infrared, ultrasound, and radio
  - Disadvantages: resolution and interference
- **Sensor-network-based positioning systems**
  - Provide the indoor localization services a simple and feasible solution.

## Research Scope

- **Timing of topology determination**
  - Pre-deployment
  - Post-deployment
- **Placement methods**
  - Random deployment approach
  - Controlled placement approach
- **Types of sensor nodes**
  - Homogeneous
  - Heterogeneous

# Research Scope



PLChiu

7

# Sensor Deployment

Pre-/Post-	R/C	Homo-/Hetero-	papers	goals	considerations
Pre-	R	Homo-	[AS03]	Min sensor density	Coverage
Pre-	R	Homo-	[LRS05]	Min sensor density	Coverage
Pre-	R	Hetero-	[MRK05]	Min cost	Two types of sensors, coverage and connectivity
Pre-	C	Homo-	[DC03] [DC102]	Min number of sensors	Coverage threshold, grid based method
Pre-	C	Homo-	[RST04]	Min number of sensors	Identifying cod, target location, robust
Pre-	C	Homo-	[[CCZ05]	Min number of sensors	Lifetime, cost
Pre-	C	Homo-	[BXX06]	Min number of sensors	Coverage, 2-connectivity
Pre-	C	Homo-	[GCB06]	Energy efficiency	Data distortion, connectivity
Pre-	C	Homo-Hetero-	[CIQ02]	Min number of sensors	Coverage, target location
Post-	R	Homo-	[SP01] [AGP04]	Energy-efficiency	K-cover
Post-	R	Hetero-	[CWL05] [DVZ06]	Max number of covers	Adjustable sensing radius, reduce radius
Post-	C	Homo-	[ZC03a]	coverage	Mobile sensors, target location
Post-	C	Hetero-	[XWH05]	Min cost	Add relay nodes, lifetime, connectivity, sensing/ relay nodes
Post	R/C	Homo-	[VVB04]	coverage	Incremental deployment
Post-	R/C	Homo-	[KGG06]	Min number of sensors	Coverage, communication efficiency
Post-	R/C	Hetero-	[IMP05]	Number of cluster	Lifetime, two-level

PLChiu

8

## Research Problems

- Difficult problems?
  - The sensor deployment problem where subject to coverage, has been transformed to the classical **set-cover/set- $k$ -cover** problem, that is **NP-complete**.
  
  - We have more intractable problems
    - Both surveillance and target positioning ability are considered

## Research Methodology

- **Mathematical models**
  - Optimization problems
  
- **Solution methods**
  - Lagrangean relaxation (LR) approach
  - Simulated annealing (SA) approach

## Research Background

- QoS in WSNs
- Coverage models
- Target location

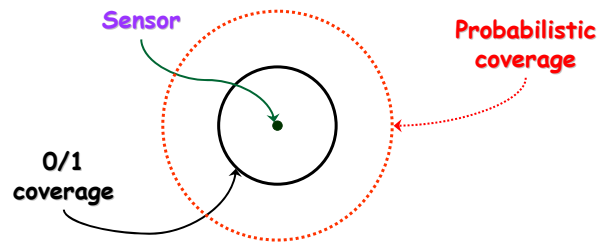
## QoS in WSNs

- In reference [WLY06], the authors define the QoS parameters for each function layer.
- Using **positioning ability** as QoS parameter
  - It belongs to application layer QoS parameter

# Coverage

## ■ Sensor coverage models

- **0/1 model**: [CIQC02], [MP03]
- Probabilistic model: [NKJ05], [DC03]



PLChiu

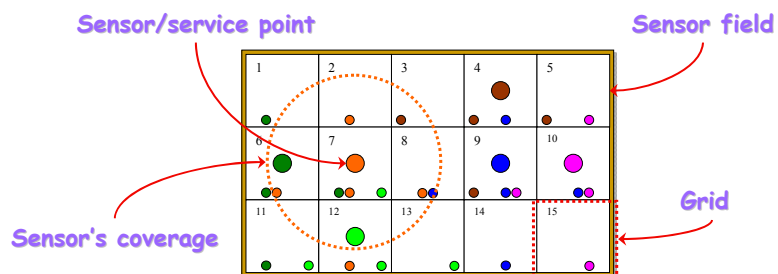
Ref: [CIQC02], [MP03], [NKJ05], [DC03], [CW04]

13

# Target Location

## ■ Grid-based sensor field

- Sensor field can be represented as a collection of two- or three-dimensional *grid points*.



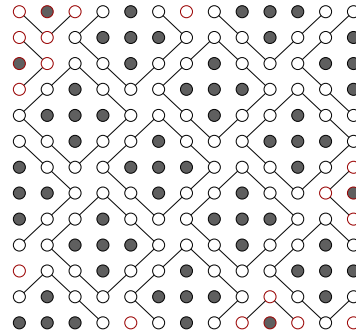
PLChiu

Ref: [RST04], [KCL98], [CIQ01], [CIQ02], [RUP03], [DC02], [LRS05], [SSS03]

14

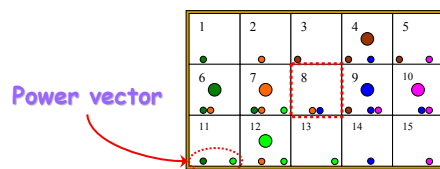
# Target Location

- Identifying code
  - The theory is adopted to design a location system with sensor network.
  - The identifying code optimization problem is NP-Complete.



# Target Location

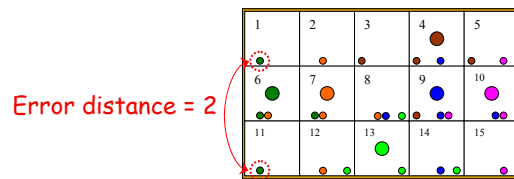
- Power vector
  - Identifying code
    - The power vector of service point 8
      - $\langle 0, 0, 1, 1, 0, 0 \rangle$  corresponding to sensors (4, 6, 7, 9, 10, 12)
- Completely covered sensor field
- Completely discriminated sensor field





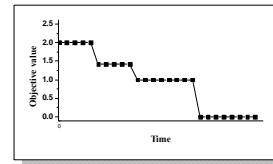
# Target Location

- Error distance
  - One of the most natural criteria to measure positioning accuracy.
- Complete discrimination
  - Maximum error distance in sensor field  $< 1$

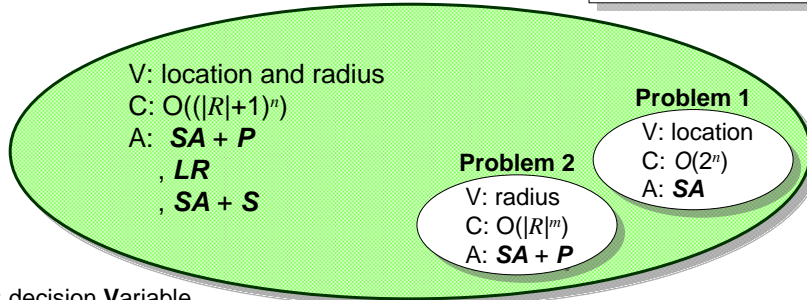


PLChiu • F.Y.S. Lin and P.L. Chiu, "A Near-optimal Sensor Placement Algorithm to Achieve Complete Coverage/Discrimination in Sensor Networks".  
 • P.L. Chiu and Frank Y.S. Lin, "A Simulated Annealing Algorithm to Support the Sensor Placement for Target Location". 17

# Problem Overview



## Problem 3



V: decision Variable  
 C: time Complexity  
 A: solution Approach  
 P: Penalty function  
 S: Surrogate function

$|R|$ : number of candidate radii  
 $n$ : number of service points  
 $m$ : number of deployed sensors,  $m \leq n$

## Research Problem 1

### ■ Sensor Placement Problem for Complete Coverage/Discrimination

- F.Y.S. Lin and **P.L. Chiu**, “A Near-optimal Sensor Placement Algorithm to Achieve Complete Coverage/Discrimination in Sensor Networks,” *IEEE Communications Letters*, vol. 9, no. 1, January 2005, pp. 43-45.
- **P.L. Chiu** and Frank Y.S. Lin, “A Simulated Annealing Algorithm to Support the Sensor Placement for Target Location,” in *Proc. IEEE CCECE 2004*, pp. 867-870.

## Problem Description

- **Objective**
  - Complete discrimination/minimizing the maximum error distance
- **Given**
  - Sensor field, set of service points, sensor cost, and detection range of sensor
- **Constraints**
  - Complete coverage and budget
- **Outcomes**
  - Sensors' location and power vectors
- **Solution approaches**
  - Simulated annealing method

## Mathematical Model

### Given Parameters:

$A$  : Index set of the sensors' candidate locations.

$B$  : Index set of the locations in the sensor field.

$r_k$  : Detection radius of the sensor located at  $k$ ;  $k \in A$ .

$d_{ij}$  : Euclidean distance between location  $i$  and  $j$ ;  $i, j \in B$ .

$c_k$  : The cost of the sensor allocated at location  $k$ ;  $k \in A$ .

$G$  : Budget limitation.

$K$  : A large number.

## Mathematical Model

### Decision Variables:

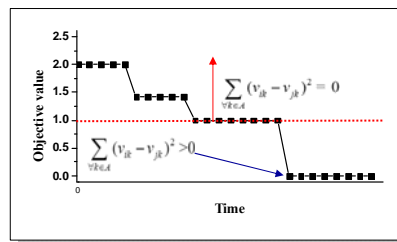
$y_k$  : 1, if a sensor is allocated at location  $k$  and 0 otherwise,  
 $k \in A$ .

$v_i = (v_{i1}, v_{i2}, \dots, v_{ik})$  : The power vector of location  $i$ , where  $v_{ik}$  is 1 if the target at location  $i$  can be detected by the sensor at location  $k$  and 0 otherwise, where  $i \in B, k \in A$ .

## Mathematical Model

### Objective function

$$Z_{IP3.1} = \min \max_{\forall k \in A} \frac{d_{ij}}{1 + K \sum_{\forall k \in A} (v_{ik} - v_{jk})^2} \quad (IP3.1)$$



## Mathematical Model

### Constraints

$$v_{ik} d_{ik} \leq y_k r_k \quad \forall k \in A, i \in B, i \neq k \quad (3.1)$$

$$\frac{d_{ik}}{r_k} > y_k - v_{ik} \quad \forall k \in A, i \in B, i \neq k \quad (3.2)$$

$$v_{ik} = y_k \quad \forall k \in A, i \in B, i \neq k \quad (3.3)$$

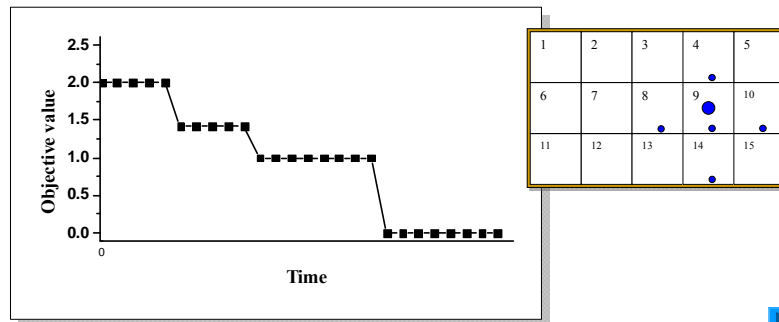
$$\sum_{\forall k \in A} c_k y_k \leq G \quad (3.4)$$

$$\sum_{\forall k \in A} v_{ik} \geq 1 \quad \forall i \in B \quad (3.5)$$

$$v_{ik}, y_k = 0 \text{ or } 1 \quad \forall k \in A, i \in B. \quad (3.6)$$

## Solution Space and Structure

- Solution space:  $O(2^n)$
- Solution structure:
  - Example: 3x5 sensor field, radius=1.



PLChiu

Sensor Placement Problem for Complete Coverage/Discrimination

25

## Simulated Annealing Based Algorithm

Algorithm 2.1: The skeleton of the SA heuristic

1. Select an energy function  $E(x)$ ;
2. Select an initial temperature  $T_0 > 0$ , and  $T = T_0$ ;
3. Select initial number of repetitions on initial temperature,  $r(T_0)$ ;
4. Set repetition counter  $t = 0$ ;
5. **Repeat**
6.   Set repetition counter  $n = 0$ ;
7.   **Repeat**
8.     Generate new state  $x_{i+1}$ , a neighbor of  $x_i$ ;
9.     Calculate  $\Delta E = E(x_{i+1}) - E(x_i)$ ;
10.     **If**  $\Delta E < 0$  **then**  $x_i = x_{i+1}$ ;
11.     **else if** random  $(0,1) < \exp(-\Delta E/T)$  **then**  $x_i = x_{i+1}$ ;
12.      $n = n + 1$ ;
13.   **Until**  $n = r(T)$ ;
14.    $t = t + 1$ ;
15.    $T_i = \alpha * T_{i-1}$ ,  $r(T_i) = \beta * r(T_{i-1})$ ;
16. **Until** stopping criterion,  $T_i < T_f$  is true.

$$E = \max_{v \in A} \frac{d_{ij}}{1 + K \sum (v_{ik} - v_{jk})^2}$$

**Dropping one sensor  
or  
Altering one sensor's  
position**

PLChiu

Sensor Placement Problem for Complete Coverage/Discrimination

26

## Experimental Results

- Scenario
  - Pentium IV-1.4GHz PC
  - Microsoft Window XP Pro.
  - Microsoft Visual C++ 6.0
  - Sensor radius = 1,  $c_i=1, \forall 1 \leq i \leq n$ .
  - The parameters of the cooling schedule:
    - $\alpha = 0.75, \beta = 1.3, b_0 = 5n$ , (where  $n$  is the amount of grids in the sensor field.
    - $t_0 = 0.1, t_f = t_0/30$
    - $K = 10000$

## Experiment I

Table 3.1: Comparison between exhaustive search and the SA algorithm.

Area	# of sensors		Sensor density	Area	# of sensors		Sensor density
	Opt.	SA			Opt.	SA	
3x3	4	4	44.44%	6x4	10	10	41.67%
4x3	6	6	50.00%	6x5	12	12	40.00%
4x4	7	7	43.75%	7x3	9	9	42.86%
5x3	6	6	40.00%	7x4	12	12	42.86%
5x4	8	8	40.00%	8x3	10	10	41.67%
5x5	10	10	40.00%	9x3	11	11	40.74%
6x3	8	8	44.44%	10x3	12	12	40.00%

## Experiment II

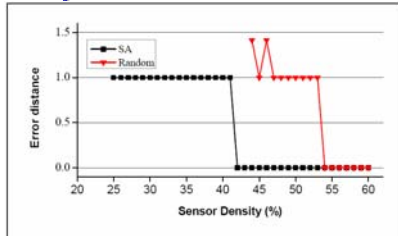


Figure 3.2: Error distance vs. sensor density. (10x10, R=1)

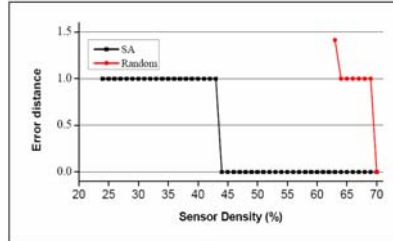


Figure 3.3: Error distance vs. sensor density. (30x30, R=1)

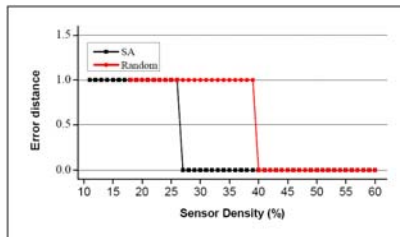


Figure 3.4: Error distance vs. sensor density. (10x10, R=2)

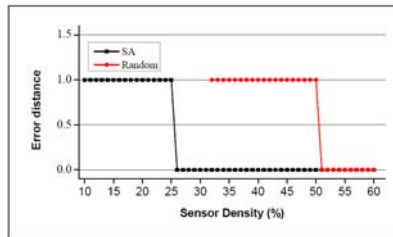


Figure 3.5: Error distance vs. sensor density. (30x30, R=2)

## Research Problem 2

- Random Deployment Problem using Adjustable Sensing Radius

## Problem Description

- **Objective**
  - Complete discrimination/minimizing the maximum error distance
- **Given**
  - Sensor field, set of service points, number of sensor, location of each sensor, and set of candidate detection radius
- **Constraints**
  - Complete coverage
- **Outcomes**
  - Sensors' radius and power vectors
- **Solution approach**
  - Simulated annealing method

## Experiment

- **Step 1: the initial configurations**
  - Random approach
  - 1000 samples
- **Step 2: adjusting radius**
  - The proposed algorithm
  - 100 selected configurations with complete coverage



# Experiment

## ■ Scenario

- Pentium IV-1.4GHz PC
- Microsoft Window XP Pro.
- Microsoft Visual C++ 6.0
- Sensor field is 150 (10 by 15) service points.
- Sensor cost,  $c_i=1, \forall 1 \leq i \leq n$ .
- Set of candidate radius  $R= \{1, 2, \dots, 8\}$ .
- The parameters of the cooling schedule:
  - $\alpha = 0.7, \beta = 1.3, b_0 = 2000$ .
  - $t_0 = 1.0, t_f = t_0/2000$
  - $K = 10000$

# Results: Adjusting Radii

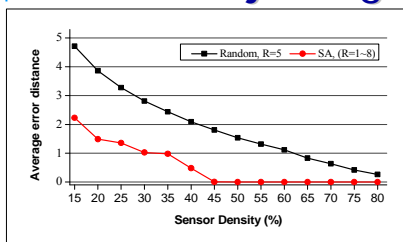


Figure 3.10: Average error distance.

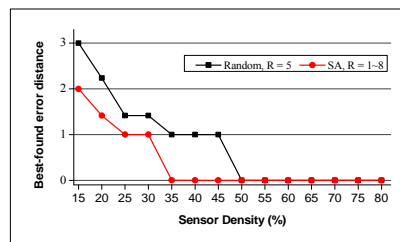


Figure 3.11: Minimum error distance.

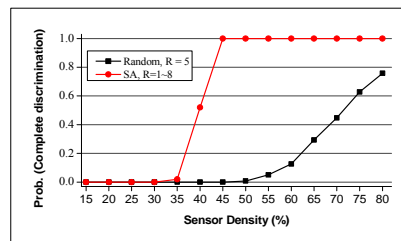


Figure 3.12: Probability for achieving complete discrimination .

## Conclusion Remarks (Problem 1 and 2)

- Mathematical model
- Simulated annealing algorithm (Algorithm 3.1)
  - Efficiently obtain high-quality solutions
  - Very effective and scalable
- Simulated annealing algorithm (Algorithm 3.2)
  - Very effective for reducing the error distance and improving the probability of complete discrimination

## Problem 3

- Sensor Placement Problem for Differentiated Quality of Positioning Services
  - **P.L. Chiu** and F.Y.S. Lin, "Sensor Placement Algorithms Supporting Differentiated Positioning Services," Submitted to *Globecom 2007 Ad-hoc and Sensor Networking Symposium*.

## Problem Description

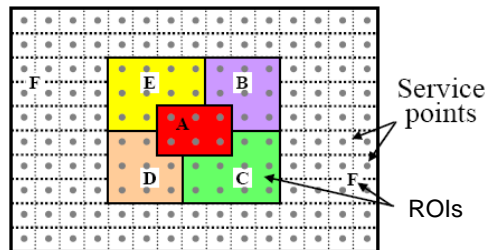
- Differentiated QoS requirements
  - QoS parameter: positioning accuracy
  - Discrimination weight
- Objective
  - Satisfying QoS requirement for all ROIs/minimizing QoS degradation
- Given
  - Sensor field, set of service points, deployment cost for each location, and set of candidate detection ranges
- Constraints
  - Budget, complete coverage, placement limitation, and service priority for each ROI (regions of interest)
- Outcomes
  - Sensors' placement, radius and power vectors
- Solution approaches
  - Lagrangean relaxation method and simulated annealing method

## Framework

- Grid based placement
- Placement limitation
- Adjustable detection radius
- ROIs (regions of interest)
- Service priority of ROI
- Types of QoS
  - Completely discriminated
  - Discriminable
  - Surveillance-only
- Goal
  - Satisfy QoS requirement for all ROIs
  - Reduce QoS degradation

## Example (museum)

- Three priority classes: high, medium, and low
  - Area A: high service priority
  - Area B, C, D, and E: medium service priority
  - Area F: low service priority



## Example (museum)

- Levels of QoS

Table 4.1: The QoS supported for ROIs in the museum example.

Level of QoS	QoS supported for ROIs
1	<ul style="list-style-type: none"> <li>● <b>Completely discriminable:</b> None.</li> <li>● <b>Discriminable:</b> ROI A.</li> <li>● <b>Surveillance-only:</b> ROIs B, C, D, E, and F.</li> </ul>
2	<ul style="list-style-type: none"> <li>● <b>Completely discriminable:</b> ROI A.</li> <li>● <b>Discriminable:</b> ROIs B, C, D, and E.</li> <li>● <b>Surveillance-only:</b> ROI F.</li> </ul>
3	<ul style="list-style-type: none"> <li>● <b>Completely discriminable:</b> ROIs A, B, C, D, and E.</li> <li>● <b>Discriminable:</b> ROI F.</li> </ul>
4	<ul style="list-style-type: none"> <li>● <b>Completely discriminable:</b> ROIs A, B, C, D, E, and F.</li> </ul>

## Mathematical Model

### Given Parameters:

- $A$  : Index set of the service points in the sensor field.
- $B$  : Index set of the sensor's candidate locations.  $B \subseteq A$ .
- $C$  : Set of the kinds of cost for sensor
- $W$  : Set of the discrimination weight
- $R$  : Set of candidate detection radiuses for sensor
- $d_{ij}$  : Euclidean distance between location  $i$  and  $j$ ;  $i, j \in A$ .
- $c_k$  : The cost of sensor located at position  $k$ ;  $k \in B$ ,  $c_k \in C$ .
- $c_{min}$  : The minimum cost of sensors.
- $G$  : The budget limitation for sensors.
- $N$  : The maximum number of sensors,  $N = G/c_{min}$ .
- $w_{ij}$  : Discrimination weight,  $i, j \in A$ ,  $w_{ij} \in W$ .
- $K$  : A larger number

## Mathematical Model

### Decision Variables:

- $y_k$  : 1, if a sensor is allocated at position  $k$ , and 0 otherwise,  $k \in B$ .
- $v_i = (v_{i1}, v_{i2}, \dots, v_{ik})$  : A power vector of location  $i$ , where  $v_{ik}$  is 1 if the target at location  $i$  can be detected by the sensor at position  $k$  and 0 otherwise,  $i \in A$ ,  $k \in B$ .
- $r_k$  : Detection radius of sensor located at  $k$ ,  $k \in B$ .

## Mathematical Model

- The objective function for the original Problem

$$Z_{IP4.1} = \min_v \max_{(i,j)} \frac{w_{ij} d_{ij}}{1 + K \sum_{\forall k \in B} (v_{ik} - v_{jk})^2} \quad (IP4.1)$$

## Mathematical Model

- Constraints

$$v_{ik} d_{ik} \leq y_k r_k \quad \forall i \in A, k \in B, i \neq k \quad (4.1)$$

$$\frac{d_{ik}}{r_k} > y_k - v_{ik} \quad \forall i \in A, k \in B, i \neq k \quad (4.2)$$

$$v_{kk} = y_k \quad \forall k \in A \cap B \quad (4.3)$$

$$\sum_{\forall k \in B} c_k y_k \leq G \quad (4.4)$$

$$\sum_{\forall k \in B} v_{ik} \geq 1 \quad \forall i \in A \quad (4.5)$$

$$\sum_{\forall k \in B} v_{ik} \leq N \quad \forall i \in A \quad (4.6)$$

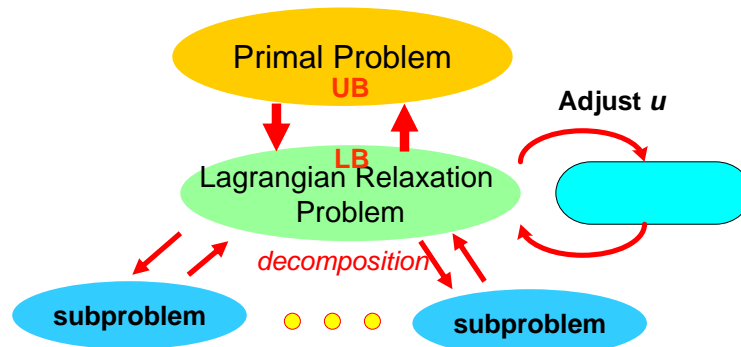
$$r_k \in R \quad \forall k \in B \quad (4.7)$$

$$v_{ik} = 0 \text{ or } 1 \quad \forall i \in A, k \in B \quad (4.8)$$

$$y_k = 0 \text{ or } 1 \quad \forall i \in A, k \in B. \quad (4.9)$$

# Lagrangian Relaxation

$$LB \leq \text{Optimal solution} \leq UB$$



# Equivalent Model

- The original model  $Z_{IP4.1} = \min_v \max_{(i,j)} \frac{w_{ij} d_{ij}}{1 + K \sum_{\forall k \in B} (v_{ik} - v_{jk})^2}$

is hard to solve directly.

- Define  $S = \max_{\forall k \in A} \frac{w_{ij} d_{ij}}{1 + K \sum_{\forall k \in A} (v_{ik} - v_{jk})^2}$

$$\text{then } S \geq \frac{w_{ij} d_{ij}}{1 + K \sum_{\forall k \in A} (v_{ik} - v_{jk})^2} \text{ and } \frac{1}{S} \leq \frac{1 + K \sum_{\forall k \in A} (v_{ik} - v_{jk})^2}{w_{ij} d_{ij}}$$

## Equivalent Model

$$Z_{IP4.2} = \min_v S \quad (IP4.2)$$

subject to:

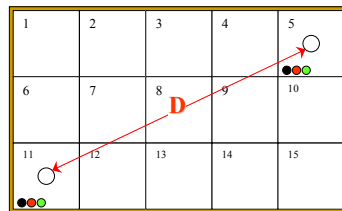
$$\frac{1}{S} \leq \frac{1 + K \sum_{\forall k \in B} (v_{ik} - v_{jk})^2}{w_{ij} d_{ij}} \quad \forall i, j \in A, i \neq j \quad (4.11)$$

$$\underline{S} \leq S \leq \bar{S} \quad \begin{array}{l} \text{is lower bound of } S; \\ \text{is upper bound of } S. \end{array} \quad (4.12)$$

and (4.1)-(4.9)

## Lemma 4.3

■ *Upper bound*  $\bar{S} = \max_{\forall i, j \in A} \frac{w_{ij} d_{ij}}{1 + K \sum_{\forall k \in A} (v_{ik} - v_{jk})^2} \leq w_h D$



Hamming distance = 0



## Lemma 4.4

■ **Lower bound**  $\underline{S} \leq \max_{\forall i, j \in A} \frac{w_{ij} d_{ij}}{1 + K \sum_{\forall k \in A} (v_{ik} - v_{jk})^2}$ ,  $\underline{S} = \frac{w_{\ell}}{1 + KD_h}$

, where  $D_h = \max\{2n, N\}$ .

• Case  $N \geq 2n$ : the maximum Hamming distance is  $2n$ .

E.g.,

$$\begin{aligned} \text{Power vector A} &= \langle v_0, v_1, \dots, v_{n-1}, 0, \dots, 0, \underbrace{0, \dots, 0}_{n \text{ elements}} \rangle \\ \text{Power vector B} &= \langle \underbrace{0, 0, \dots, 0}_{n \text{ elements}}, 0, \dots, 0, \underbrace{v_{N-1-n-1}, \dots, v_{N-1}}_{n \text{ elements}} \rangle \end{aligned}$$

• Case  $N < 2n$ : the maximum Hamming distance is  $N$ .

E.g.,

$$\begin{aligned} \text{Power vector A} &= \langle v_0, v_1, \dots, v_{n-1}, 0, \dots, 0 \rangle \\ \text{Power vector B} &= \langle \underbrace{0, 0, \dots, 0}_{n \text{ elements}}, \underbrace{v_{N-1-n-1}, \dots, v_{N-1}}_{n \text{ elements}} \rangle \end{aligned}$$

## Transformation

$$\frac{1}{S} \leq \frac{1 + K \sum_{k \in A} (v_{ik} - v_{jk})^2}{w_{ij} d_{ij}} \quad (4.11)$$

$$\frac{1}{S} \leq \frac{1 + \sum_{k \in A} (v_{ik} - 2v_{ik}v_{jk} + v_{jk})}{w_{ij} d_{ij}}$$

■ The cutting plane method is applied to transform Constraint (4.11) to a linear Constraint (4.13).

■ An auxiliary variable,  $t_{ijk}$ , is introduced, where

$$t_{ijk} = v_{ik}v_{jk}.$$

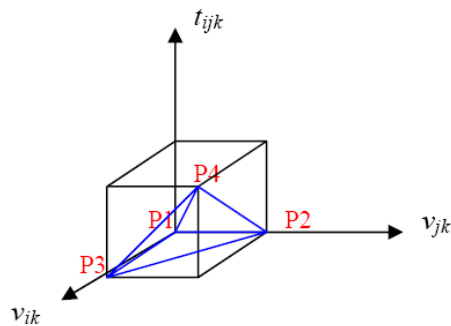
## Transformation

- The truth table for variables  $v_{ik}$ ,  $v_{jk}$ , and  $t_{ijk}$ .

$v_{ik}$	$v_{jk}$	$t_{ijk}$
0	0	0
0	1	0
1	0	0
1	1	1

## Transformation

- The possible values for the three variables only exist in four integer vertices,  $p_1$ ,  $p_2$ ,  $p_3$ , and  $p_4$ , of the polyhedron.



## Transformation

- The *four planes* constructing the polyhedron are presented as following.
- Relationship between  $v_{ik}$ ,  $v_{jk}$ , and  $t_{ijk}$ .

$$\begin{array}{ll}
 v_{ik} - t_{ijk} \geq 0 & \forall i, j \in A, i \neq j, k \in B \\
 v_{jk} - t_{ijk} \geq 0 & \forall i, j \in A, i \neq j, k \in B \\
 v_{ik} + v_{jk} - t_{ijk} \leq 1 & \forall i, j \in A, i \neq j, k \in B \\
 t_{ijk} \geq 0 & \forall i, j \in A, i \neq j, k \in B
 \end{array}$$

## Transformation

*Let*  $t_{ijk} = v_{ik}v_{jk}$

*then*  $(v_{ik} - v_{jk})^2 = v_{ik} + v_{jk} - 2t_{ijk}$

Rewrite Constraint (4.11)

$$\frac{1}{S} \leq \frac{1 + K \sum_{k \in B} (v_{ik} - v_{jk})^2}{w_{ij} d_{ij}}$$

*to* 
$$\frac{1}{S} \leq \frac{1 + K \sum_{k \in B} (v_{ik} + v_{jk} - 2t_{ijk})}{w_{ij} d_{ij}}$$

$$Z_{IP4.3} = \min_v S \quad (IP4.3)$$

Subject to:

$$\frac{1}{S} \leq \frac{1 + K \sum_{\forall k \in B} (v_{ik} + v_{jk} - 2t_{ijk})}{w_{ij} d_{ij}} \quad \forall i, j \in A, i \neq j \quad (4.12)$$

$$\underline{S} \leq S \leq \bar{S} \quad \begin{array}{l} \underline{S} \text{ is lower bound of } S; \\ \bar{S} \text{ is upper bound of } S. \end{array} \quad (4.1)$$

$$v_{ik} d_{ik} \leq y_k r_k \quad \forall i \in A, k \in B, i \neq k \quad (4.2)$$

$$v_{kk} = y_k \quad \forall k \in A \cap B \quad (4.3)$$

$$\sum_{\forall k \in B} c_k y_k \leq G \quad (4.4)$$

$$\sum_{\forall k \in B} v_{ik} \geq 1 \quad \forall i \in A \quad (4.5)$$

$$\sum_{\forall k \in B} v_{ik} \leq N \quad \forall i \in A \quad (4.6)$$

$$r_k \in R \quad \forall k \in B \quad (4.7)$$

$$v_{ik} = 0 \text{ or } 1 \quad \forall i \in A, k \in B \quad (4.8)$$

$$y_k = 0 \text{ or } 1 \quad \forall i \in A, k \in B \quad (4.9)$$

$$\sum_{\forall k \in B} t_{ijk} \leq N \quad \forall i, j \in A, i \neq j \quad (4.13)$$

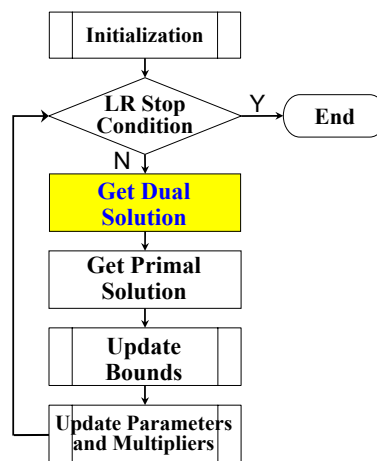
$$v_{ik} - t_{ijk} \geq 0 \quad \forall i, j \in A, i \neq j, k \in B \quad (4.14)$$

$$v_{jk} - t_{ijk} \geq 0 \quad \forall i, j \in A, i \neq j, k \in B \quad (4.15)$$

$$v_{ik} + v_{jk} - t_{ijk} \leq 1 \quad \forall i, j \in A, i \neq j, k \in B \quad (4.16)$$

$$t_{ijk} = 0 \text{ or } 1 \quad \forall i, j \in A, i \neq j, k \in B \quad (4.17)$$

## Procedure of the LR Approach



## Dual Problem

- Transform the primal problem (IP4.3) into the LR problem (LR 4.2) by relax constraints (4.1)-(4.3), (4.11), (4.14)-(4.16).
- The multipliers  $\{u_{ij}^1\}, \{u_{ik}^2\}, \{u_{jk}^3\}, \{u_{kk}^4\}, \{u_{jk}^5\}, \{u_{ijk}^6\}, \{u_{ijk}^7\}$
- Decomposition
  - Subproblems (SUB 4.1)-(SUB 4.4)
- Optimally solving each subproblem

## Subproblem 1 (for $S$ )

$$Z_{SUB4.1}(u^1) = \min \left( S + \sum_{\forall i \in A} \sum_{\substack{\forall j \in A \\ i \neq j}} u_{ij}^1 \frac{1}{S} \right) \quad (SUB4.1)$$

subject to:

$$\underline{S} \leq S \leq \bar{S} \quad (4.11)$$

- To optimal solve the subproblem, the right hand side of Equation (SUB4.1) will be differentiated respected to variable  $S$ .
- Let the new equation equals to zero and get the optimal solution of variable  $S$ ,

$$S_{opt} = \sqrt{\sum_{\forall i \in A} \sum_{\substack{\forall j \in A \\ i \neq j}} u_{ij}^1}$$

## Subproblem 1 (for $S$ )

- If  $\underline{s} \leq s_{opt} \leq \bar{s}$  then let  $Z_{SUB4.1} = 2 \sqrt{\sum_{\substack{\forall i \in A \\ \forall j \in A \\ i \neq j}} u_{ij}^1}$
- Otherwise,  $\underline{s}$  and  $\bar{s}$  are substituted for  $S$ .
- We can get optimal solution such that

$$Z_{SUB4.1}(u^1) = \min \left\{ \left( \bar{s} + \sum_{\substack{\forall i \in A \\ \forall j \in A \\ i \neq j}} \frac{u_{ij}^1}{\bar{s}} \right), \left( \underline{s} + \sum_{\substack{\forall i \in A \\ \forall j \in A \\ i \neq j}} \frac{u_{ij}^1}{\underline{s}} \right) \right\}$$

## Subproblem 2 (for $y_k, r_k$ )

$$Z_{SUB4.2}(u^2, u^3, u^4) = \min \sum_{\forall k \in B} \left( \sum_{\substack{\forall i \in A \\ i \neq k}} \left( -u_{ik}^2 r_k + u_{ik}^3 \right) y_k - u_{ik}^3 \frac{d_{ik}}{r_k} - \sum_{\forall k \in B} u_{kk}^4 y_k \right) \quad (SUB4.2a)$$

subject to:

$$\sum_{\forall k \in B} c_k y_k \leq G \quad (4.4)$$

$$r_k \in R \quad \forall k \in B \quad (4.7)$$

$$y_k = 0 \text{ or } 1 \quad \forall i \in A, k \in B. \quad (4.9)$$

## Subproblem 3 (for $v_{ik}$ )

$$\begin{aligned}
 Z_{SUB4.3}(u^1, u^2, u^3, u^4, u^5, u^6, u^7) = \min & \left( -K \sum_{i \in A} \sum_{\substack{j \in A \\ j \neq i}} u_{ij}^1 \frac{\sum_{k \in B} v_{ik}}{w_{ij} d_{ij}} \right. \\
 & -K \sum_{i \in A} \sum_{\substack{j \in A \\ j \neq i}} u_{ij}^1 \frac{\sum_{k \in B} v_{jk}}{w_{ij} d_{ij}} + \sum_{i \in A} \sum_{\substack{k \in B \\ j \neq k}} u_{ik}^2 d_{ik} v_{ik} - \sum_{i \in A} \sum_{\substack{k \in B \\ j \neq k}} u_{jk}^3 v_{ik} + \sum_{k \in K} \sum_{k \in B} u_{kk}^4 v_{ik} \\
 & - \sum_{i \in A} \sum_{\substack{j \in A \\ j \neq i}} \sum_{k \in B} u_{jk}^5 v_{ik} - \sum_{i \in A} \sum_{\substack{j \in A \\ j \neq i}} \sum_{k \in B} u_{jk}^6 v_{jk} + \sum_{i \in A} \sum_{\substack{j \in A \\ j \neq i}} \sum_{k \in B} u_{jk}^7 v_{ik} \\
 & \left. + \sum_{i \in A} \sum_{\substack{j \in A \\ j \neq i}} \sum_{k \in B} u_{jk}^7 v_{jk} \right) \quad (SUB4.3)
 \end{aligned}$$

- For each term with the variable  $v_{jk}$ , index  $j$  substitutes for  $i$  contrariwise.

## Subproblem 3 (for $v_{ik}$ )

$$\begin{aligned}
 Z_{SUB4.3}(u^1, u^2, u^3, u^4, u^5, u^6, u^7) = \min & \sum_{i \in A} \left\{ \sum_{\substack{k \in B \\ j \in A \\ j \neq i}} \left( \frac{-Ku_{ij}^1}{w_{ij} d_{ij}} - \frac{Ku_{ji}^1}{w_{ji} d_{ji}} - u_{ijk}^5 - u_{jik}^6 + u_{ijk}^7 + u_{jik}^7 \right) \right. \\
 & \left. + \sum_{\substack{k \in B \\ j \neq k}} (u_{ik}^2 d_{ik} - u_{ik}^3) + \sum_{\substack{k \in B \\ j=k}} u_{ik}^4 \right\} v_{ik} \quad (SUB4.3a)
 \end{aligned}$$

subject to:

$$\sum_{k \in B} v_{ik} \geq 1 \quad \forall i \in A \quad (4.5)$$

$$\sum_{k \in B} v_{ik} \leq N \quad \forall i \in A \quad (4.6)$$

$$v_{ik} = 0 \text{ or } 1 \quad \forall i \in A, k \in B \quad (4.8)$$

## Subproblem 4 (for $t_{ijk}$ )

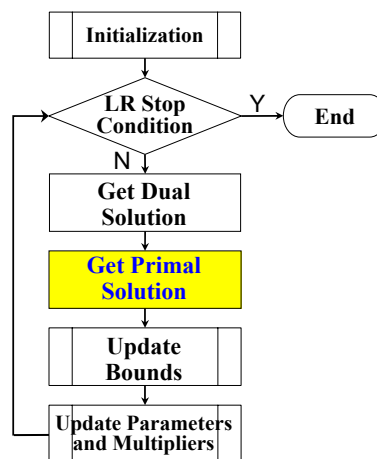
$$Z_{SUB4.4}(u^1, u^5, u^6, u^7) = \min \sum_{\substack{\forall i \in A \\ j \neq i}} \sum_{\forall j \in A} \sum_{\forall k \in B} \left( \frac{2Ku_{ij}^1}{w_{ij}d_{ij}} + u_{ijk}^5 + u_{ijk}^6 - u_{ijk}^7 \right) t_{ijk} \quad (SUB4.4a)$$

subject to:

$$\sum_{\forall k \in B} t_{ijk} \leq N \quad \forall i, j \in A, i \neq j \quad (4.13)$$

$$t_{ijk} = 0 \text{ or } 1 \quad \forall i, j \in A, i \neq j, k \in B \quad (4.17)$$

## Procedure of the LR Approach

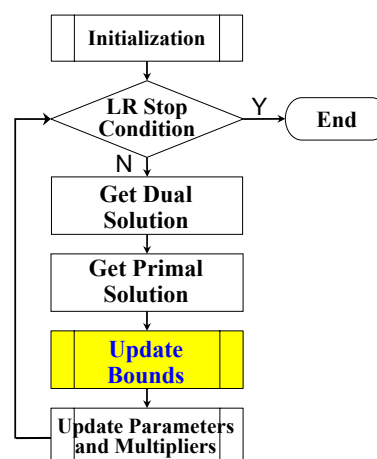




## Getting Primal Feasible

- Adjusting the dual solutions
  - Step 1: initializing the decision variables ( $y_k, v_{ik}, r_k$ )
  - Step 2: for satisfying the coverage and budget constraints, the sensor might be *added, deleted, or changed radius*.
  - Step 3: in order to improving discrimination, the sensor might be *changed position, modified radius*.

## Procedure of the LR Approach



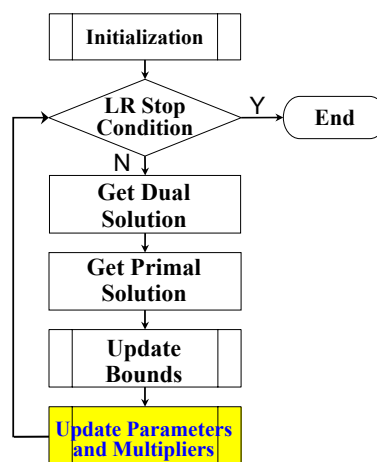
## LB

$$Z_D(u^1, u^2, u^3, u^4, u^5, u^6, u^7) = Z_{SUB4.1} + Z_{SUB4.2} + Z_{SUB4.3} + Z_{SUB4.4} + \left( - \sum_{\substack{\forall i \in A \\ i \neq j}} \sum_{\substack{\forall j \in A \\ i \neq j}} \frac{u_{ij}^1}{w_{ij} d_{ij}} - \sum_{\substack{\forall i \in A \\ i \neq j}} \sum_{\substack{\forall j \in A \\ i \neq j}} \sum_{\forall k \in B} u_{ijk}^7 \right)$$

- According the weak Lagrangean duality theorem [Fis81], [Fis85], the optimal objective value of the dual problem (LR),  $Z_D(u^1, u^2, u^3, u^4, u^5, u^6, u^7)$ , is a lower bound on primal problem (IP4.3)

$$Z_D = \max_{(u^1, u^2, u^3, u^4, u^5, u^6, u^7) \geq 0} Z_D(u^1, u^2, u^3, u^4, u^5, u^6, u^7) \quad (D1)$$

## Procedure of the LR Approach



## Update the LR Multipliers

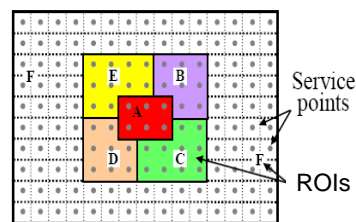
- We denote  $\pi = (u_{ij}^1, u_{ik}^2, u_{ik}^3, u_{ik}^4, u_{jk}^5, u_{jk}^6, u_{jk}^7)$  as the vector of Lagrangean multipliers with respect to relaxed constraints.
- In iteration  $m$  of subgradient optimization procedure, the multiplier  $\pi^m$  is updated by  $\pi^{m+1} = \pi^m + \zeta^m g^m$
- The step size  $\zeta^m$  is determined by

$$\zeta^m = \frac{\lambda(Z_{IP4.3}^* - Z_D(\pi^m))}{\|g^m\|^2}$$

, where  $Z_{IP4.3}^*$  represents an upper bound on the primal objective value, obtained by applying a heuristic to (IP4.3), and  $\lambda$  is a scalar satisfying  $0 \leq \lambda \leq 2$ .

## Experimental Results

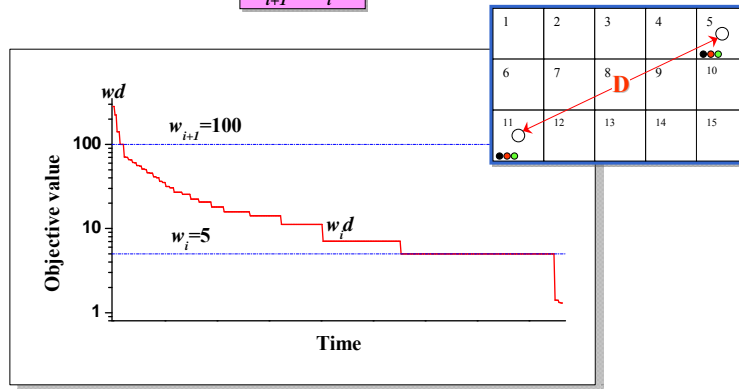
- Scenario
  - Pentium IV-3.0GHz PC
  - Microsoft Window XP Pro.
  - Microsoft Visual C++ 6.0
  - Sensor fields: 10x15
  - Sensor radius: 1 ~ 8
  - LR parameters:
    - $1 < \lambda \leq 2$
    - Improvement counter: 45
    - Amount of iteration: 1500



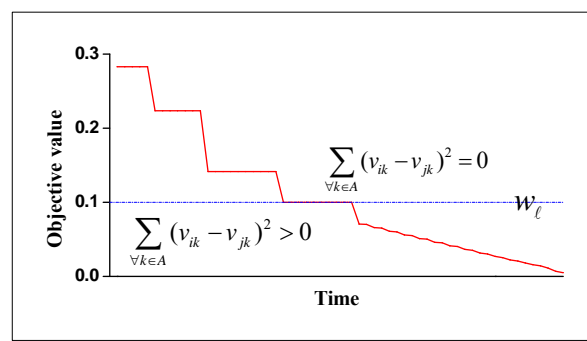
Sensor field in Exp.

## Lemma 4.1 (for $w_i$ )

$$w_{i+1} > w_i D$$



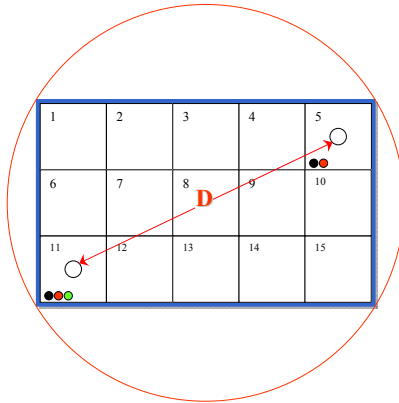
## Lemma 4.2 (for $K$ )



$$w_\ell > \max \left( \frac{w_i d_{ij}}{1 + K \sum_{\forall k \in A} (v_{ik} - v_{jk})^2} \right), \sum_{\forall k \in A} (v_{ik} - v_{jk})^2 > 0, i \neq j.$$

## Lemma 4.2 (for $K$ )

- Case 1:  $2r \geq D$

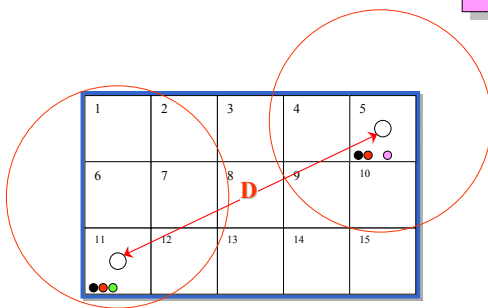


$$w_\ell > \frac{w_h D}{1 + K}$$

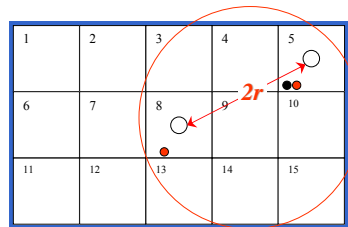
## Lemma 4.2 (for $K$ )

- Case 1:  $2r < D$

$$w_\ell > \max \left\{ \frac{w_h D}{1 + 2K}, \frac{w_h (2r)}{1 + K} \right\}$$



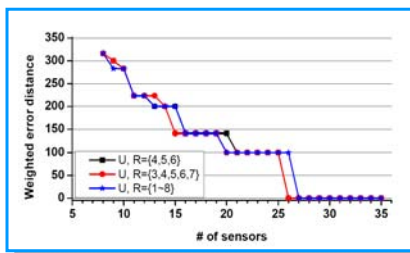
$$(1) \frac{w_h D}{1 + 2K}$$



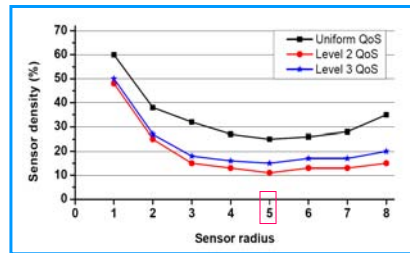
$$(2) \frac{w_h (2r)}{1 + K}$$

## Experimental Results (LR, Exp. I)

- Sensor density vs. radius
  - Sensors density depends strongly on sensor radius



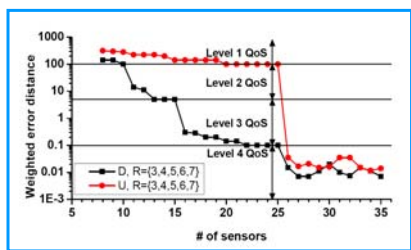
Adjustable radius, Uniform QoS



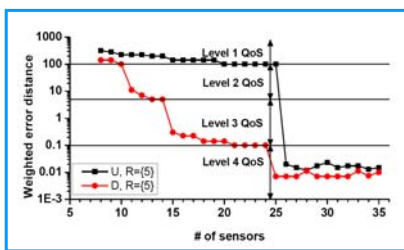
Fixed radius

## Experiment Results (LR, Exp. I)

- Uniform vs. differentiated QoS request



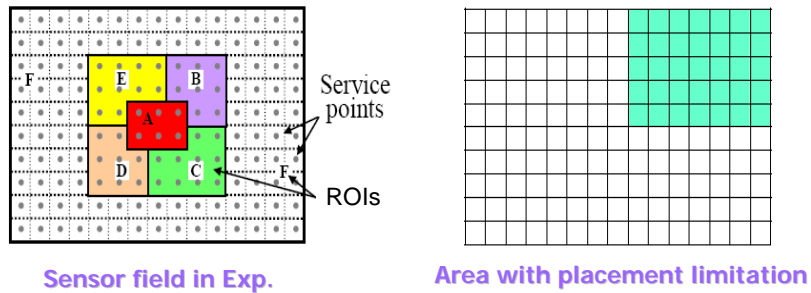
Adjustable radius



Fixed radius

## Experimental Results (LR, Exp. II)

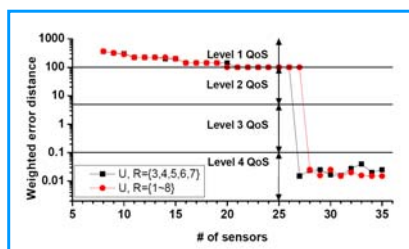
### ■ Placement Limitation



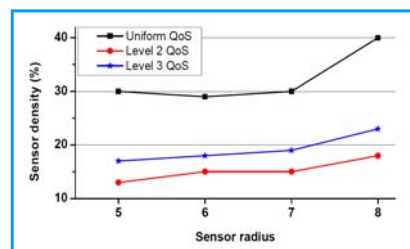
## Experimental Results (LR, Exp. II)

### ■ Sensor density vs. radius

- Sensors density depends strongly on sensor radius



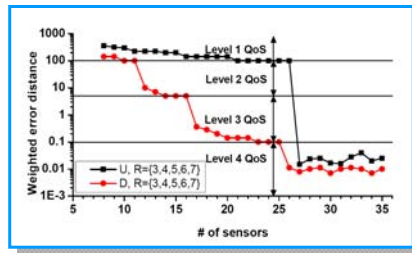
Adjustable radius, Uniform QoS



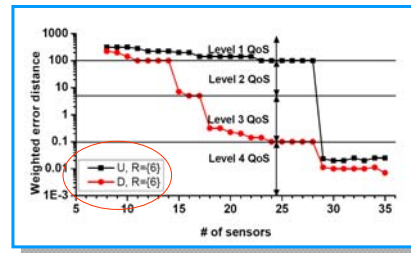
Fixed radius

## Experiment Results (LR, Exp. II)

### Uniform vs. differentiated QoS request



Adjustable radius

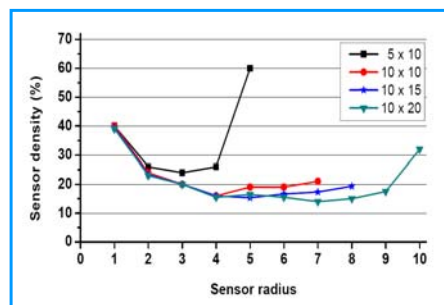


Fixed radius

## Experimental Results (LR, Exp. III)

### Scalability - solution quality

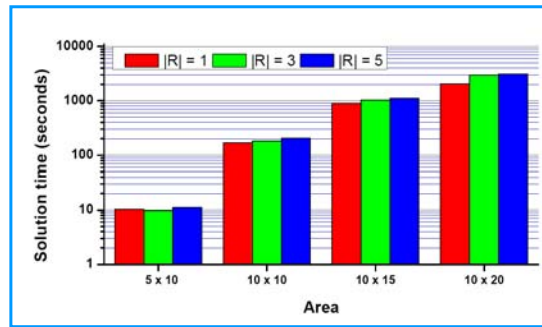
- Fixed radius
- Solution space:  $O(2^{50}) \sim O(2^{200})$





## Experimental Results (LR, Exp. III)

- Scalability – solution time
  - Solution space:  $O(2^{25}) \sim O(6^{200})$



## Simulated Annealing Approach (SA\_1)

- Penalizing objective function  $Z_{IP4.1}$  in algorithm SA\_1.
- Energy function

$$E = \max_{(i,j)} \frac{w_{ij} d_{ij}}{1 + K \sum_{\tilde{k} \in B} (v_{ik} - v_{jk})^2} \left[ 1 + p \left( g + \sum_{i \in \mathcal{A}} h_i \right) \right]$$

Penalty

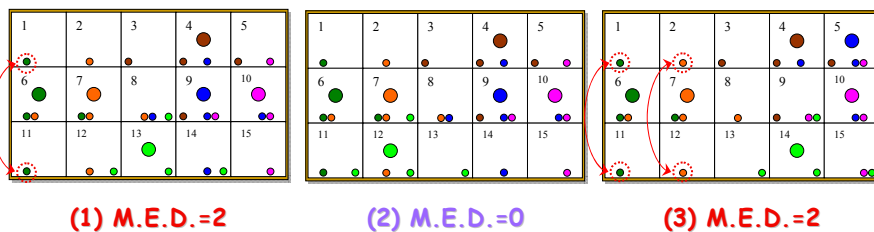
- $g$ : indicates the over budget
- $\sum h$ : number of uncovered service points
- $p$ : constant

## Simulated Annealing Approach (SA\_1)

- Actions for generating a new solution configuration
  - Adding one sensor
  - Dropping one sensor
  - Changing one sensor's location
  - Increasing one sensor's radius
  - Reducing one sensor's radius.

## Simulated Annealing Approach (SA\_2)

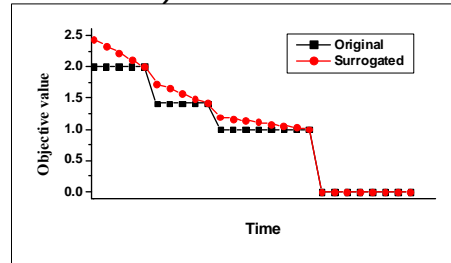
- Max. Error Distance (M.E.D.)



## Simulated Annealing Approach (SA\_2)

- Energy function (Surrogate function)

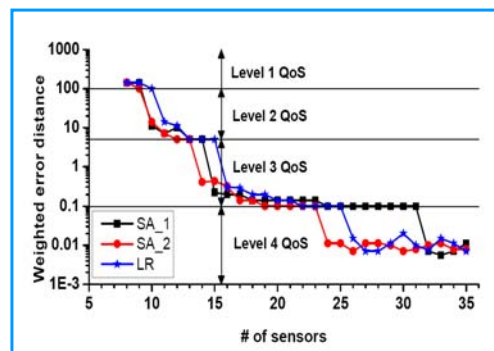
$$E_s = E * \left( 1 + \frac{p'}{|A|} \sum_{i \in A} m_i \right)$$



- $E$ : Energy function in SA\_1
- $\sum m_i / |A|$ : ratio of the service points with the maximum weighted error distance
- $p'$ : constant

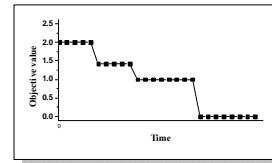
## Experimental Results (Comp.)

- Solution quality Comparison

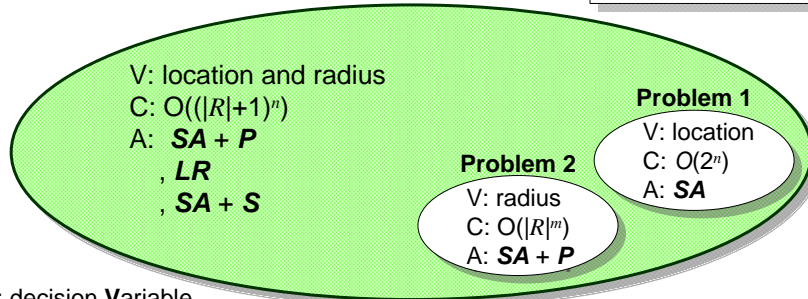


Area: 10x15, differentiated QoS, R={3, 4, 5, 6, 7}

## Discussion



### Problem 3



V: decision Variable  
 C: time Complexity  
 A: solution Approach  
**P**: Penalty function  
**S**: Surrogate function

$|R|$ : number of candidate radii  
 $n$ : number of service points  
 $m$ : number of deployed sensors,  $m \leq n$

PLChiu

87

## Problem 4

### ■ Energy-Efficient Sensor Networks Design

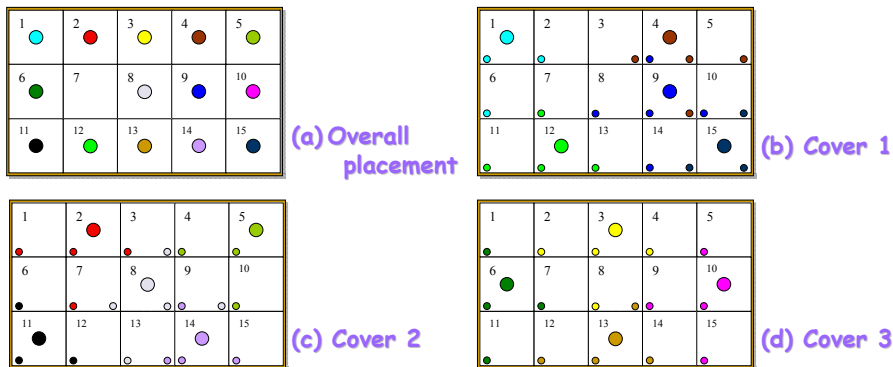
- F.Y.S. Lin and **P.L. Chiu**, “Energy-Efficient Sensor Network Design Subject to Complete Coverage and Discrimination Constraints,” in *Second Annual IEEE Communications Society Conference on Sensor and Ad Hoc Communications and Networks (IEEE SECON)*, 26-29 September 2005, pp. 586 - 593. (Acceptance ratio=55/202=27.2%)
- F.Y.S. Lin and **P.L. Chiu**, “A Simulated Annealing Algorithm for Energy Efficient Sensor Network Design,” in *3rd International Symposium on Modeling and Optimization in Mobile, Ad Hoc, and Wireless Networks (WiOpt)* April 3-7, 2005. (Acceptance ratio=44/141=31%)

PLChiu

88

# Energy-Efficiency Approach

- A grid-based sensor field with 3 covers



PLChiu

Energy-Efficient Sensor Networks Design

89

# Problem Description

- Objective
  - Minimizing the deployment cost
- Given
  - Sensor field, set of service points, sensor cost, detection radius, and number of covers
- Constraints
  - Complete coverage for each cover and complete discrimination for overall sensor network
- Outcomes
  - Sensors' placement, members of each cover, and power vectors
- Solution approach
  - Lagrangean relaxation method and simulated annealing method

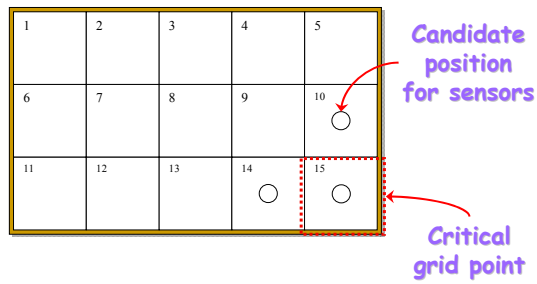
PLChiu

• F.Y.S. Lin and **P.L. Chiu**, "A Simulated Annealing Algorithm for Energy Efficient Sensor Network Design".  
 • F.Y.S. Lin and **P.L. Chiu**, "Energy-Efficient Sensor Network Design Subject to Complete Coverage and Discrimination Constraints".

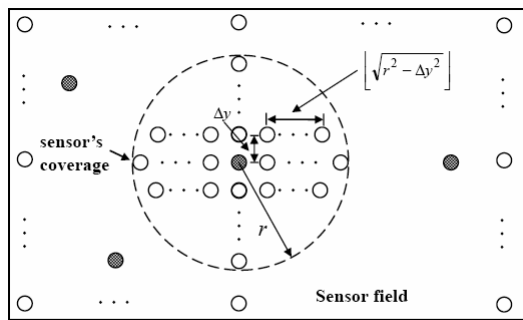
90

## Upper Bound on Number of Covers

- Critical grid point
  - Grids 1, 5, 11, and 15
  - Upper bound: 3



## Upper Bound on Number of Covers



● : sensor.      ○ : grid point.

The theoretic upper bound:

$$U_r = 2r + 1 + \sum_{\Delta y=1}^r \left\lfloor \sqrt{r^2 - \Delta y^2} \right\rfloor$$

## Mathematical Model

### Given Parameters:

- $A$  : The set of the index for candidate locations where sensor can be allocated.
- $B$  : The set of the index for service points that can be covered and located by the sensor network.
- $K$  : The number of covers required for the sensor network.
- $a_{ij}$  : Indicator function which is 1 if service point  $i$  can be covered by sensor  $j$  and 0 otherwise.
- $c_j$  : Cost function of sensor  $j$ .

## Mathematical Model

### Decision Variables:

- $x_{jk}$  : 1 if sensor  $j$  is allocated on cover  $k$  of the sensor network.
- $y_j$  : Sensor allocation decision variable which is 1 if sensor  $j$  is allocated in the sensor network.

## Mathematical Model

Problem (IP):

$$Z_{IP5.1} = \min \sum_{j \in A} \sum_{k \in K} c_j x_{jk} \quad (IP5.1)$$

Subject to:

$$\sum_{j \in A} a_{ij} x_{jk} \geq 1, \forall i \in B, k \in K \quad (5.1)$$

$$\sum_{k \in K} x_{jk} \leq 1, \forall j \in A \quad (5.2)$$

$$y_j = \sum_{k \in K} x_{jk}, \forall j \in A \quad (5.3)$$

$$\sum_{j \in A} (a_{ij} - a_{lj})^2 y_j \geq 1, \forall i, l \in B, i \neq l \quad (5.4)$$

$$x_{jk}, y_j = 0 \text{ or } 1, \forall j \in A, k \in K \quad (5.5)$$

- Solution space:  $O((K+I)^m)$

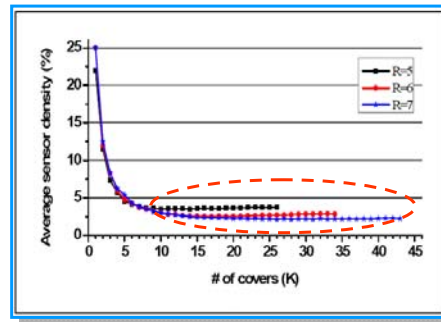
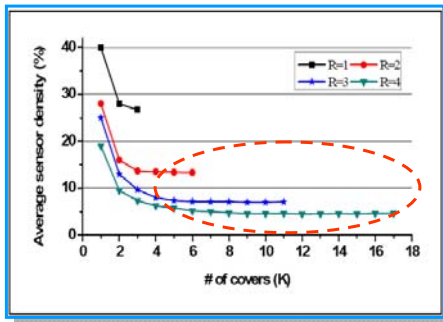
## Experimental Results (LR)

- Scenario
  - Pentium IV-1.4GHz PC
  - Microsoft Window XP Pro.
  - Microsoft Visual C++ 6.0
  - Sensor fields: 10x10 service points
  - Sensor radius: 1 ~ 7
  - LR parameters:
    - $1 < \lambda \leq 2$
    - Improvement counter: 35
    - Amount of iteration: 1000



## Experiment Results (LR)

- Average deployed density



## Experimental Results (LR)

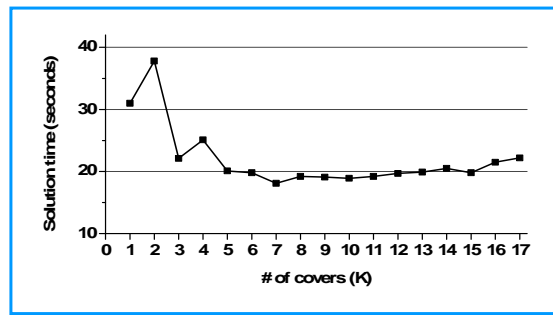
- The duplicated deployment vs. the proposed approach
- Deployment cost vs. lifetime extension

Radius	The duplicate deployment		The proposed approach	
	#Duplication	Increased cost	#Cover	Increased cost
1	3	3	3	2.00
2	6	6	6	2.71
3	11	11	11	3.04
4	17	17	17	4.16
5	26	26	26	4.41
6	34	34	34	3.88
7	43	43	43	3.88

## Experimental Results (LR)

### ■ Scalability

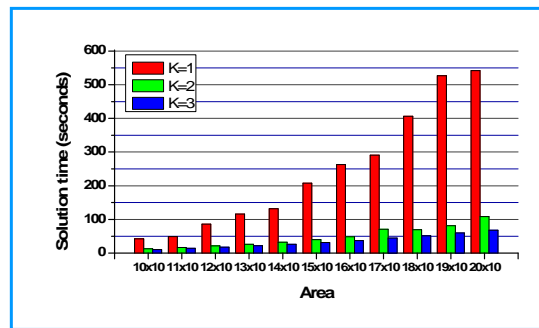
- Solution time vs. # of covers,  $R=4$
- Solution space:  $O(2^{100}) \sim O(18^{100})$



## Experimental Results (LR)

### ■ Scalability in terms of number of service points

- $r=1$
- Solution space:  $O(2^{100}) \sim O(4^{200})$



## Simulated Annealing Based Algorithm

$$E = (1 + p \sum_{\forall k \in K} \sum_{\forall i \in B} g_{ik}) (1 + p^2 (1 - d_{\min})) \sum_{\forall j \in A} \sum_{\forall k \in K} c_j x_{jk}$$

Algorithm 2.1: The skeleton of the SA heuristic.

1. Select an energy function  $E(x)$ ;
2. Select an initial temperature  $T_0 > 0$ , and  $T = T_0$ ;
3. Select initial number of repetitions on initial temperature,  $r(T_0)$ ;
4. Set repetition counter  $t = 0$ ;
5. **Repeat**
6.   Set repetition counter  $n = 0$ ;
7.   **Repeat**
8.     Generate new state  $x_{i+1}$ , a neighbor of  $x_i$ ;
9.     Calculate  $\Delta E = E(x_{i+1}) - E(x_i)$ ;
10.     **If**  $\Delta E < 0$  **then**  $x_i = x_{i+1}$ ;
11.     **else if**  $\text{random}(0,1) < \exp(-\Delta E/T)$  **then**  $x_i = x_{i+1}$ ;
12.      $n = n + 1$ ;
13.     **Until**  $n = r(T)$ ;
14.      $t = t + 1$ ;
15.      $T_i = \alpha * T_{i-1}$ ,  $r(T_i) = \beta * r(T_{i-1})$ ;
16. **Until** stopping criterion,  $T_i < T_f$  is true.

Penalty

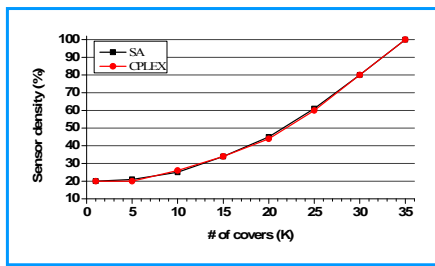
• Remove or add one sensor away/to a cover, or  
• Exchange two sensors' status

## Experimental Results (SA)

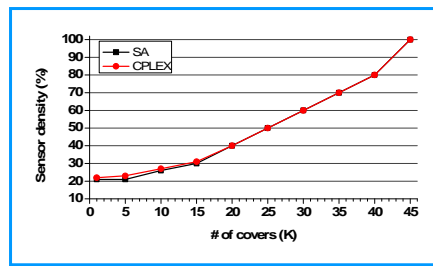
- Scenario
  - Pentium IV-1.4GHz PC
  - Microsoft Window XP Pro.
  - Microsoft Visual C++ 6.0
  - Sensor fields: 10x10 service points
  - Sensor radius: 1 ~ 7
- Benchmark: ILOG CPLEX 9.0

## Experiment Results (SA vs. CPLEX)

- Required sensor density vs. # of covers



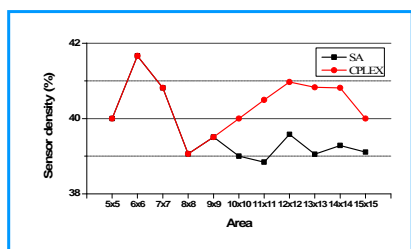
R=6



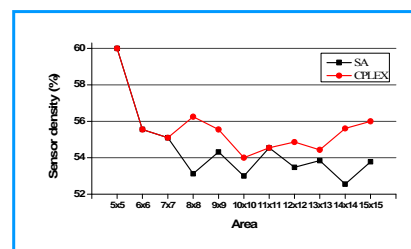
R=7

## Experimental Results (SA vs. CPLEX)

- Scalability – sensor density, R=1
  - Max. solution time for CPLEX: **10000 sec.**



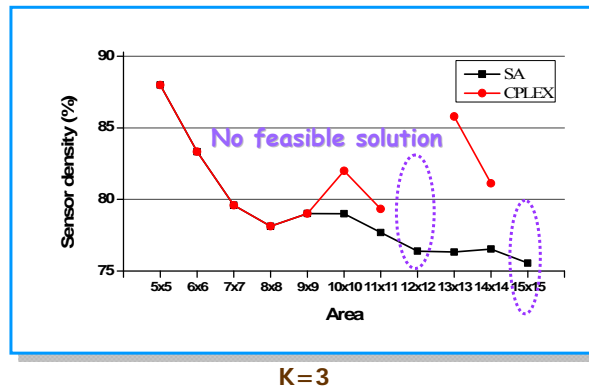
K=1



K=2

## Experimental Results (SA vs. CPLEX)

- Scalability – sensor density,  $R=1$



## Experimental Results (LR vs. SA)

- Lifetime extension

Radius (r)	1	2	3	4	5	6	7
$U_r$ (theoretic)	3	6	11	17	26	35	45
$U_r$ (SA)	3	6	11	17	26	35	45
$U_r$ (LR)	3	6	11	17	26	34	43

## Experimental Results (LR & SA)

- Solution time

Solution time (Seconds)	Sensor radius						
	1	2	3	4	5	6	7
Average (SA)	94	267	227	265	298	478	587
Average (LR)	43	61	85	38	25	91	51

## Concluding Remarks

- Energy-efficient sensor placement for surveillance and target location services
- Mathematical model
- Lagrangean relaxation method
- Simulated annealing method
- Prolonging the lifetime up to theoretical upper bound

## Concluding Remarks

- Solution time
  - The proposed LR approach outperforms the SA approach.
  - Both approaches outperform CPLEX.
- Scalability
  - Both the proposed LR and SA approaches are scalable in terms of solution quality and solution time, but CPLEX isn't.

## Contributions

- **Research scope** for sensor deployment research
- We first introduced the **positioning ability** as the QoS parameter in WSNs from application perspective, as well as proposed the error distance to measure the positioning ability.
- A **generic framework** for sensor placement problem has been proposed to support the differentiated QoS and prioritized service in WSNs.
- **Energy-efficient** sensor deployment algorithms for target positioning
- The **mathematical optimization models** were proposed to define the problems strictly.

## Contributions

- These optimization problems, since the nonlinear and combinatorial natures, are hard to solve by traditional mathematical programming methods directly. Base on LR and SA methods, we successfully developed many **heuristics** to solve these optimization problems.
- For sensor network builders and researchers, the results and algorithms proposed in this dissertation are useful.
- More relative problems can be solved by proposed methods with minor modifications.

## Future Work

- An energy-efficient deployment approach with mobile sensors
  - Model 1: for positioning
  - Model 2: for surveillance-only
- Survivable sensor networks design



## Publication List

### Journal papers:

- F.Y.S. Lin and **P.L. Chiu**, "A Near-optimal Sensor Placement Algorithm to Achieve Complete Coverage/Discrimination in Sensor Networks," *IEEE Communications Letters*, vol. 9, no. 1, January 2005, pp. 43-45.

### Conference papers:

- **P.L. Chiu** and Frank Y.S. Lin, "A Simulated Annealing Algorithm to Support the Sensor Placement for Target Location," *Proc. IEEE CCECE 2004*, pp. 867-870.
- F.Y.S. Lin and **P.L. Chiu**, "A Simulated Annealing Algorithm for Energy Efficient Sensor Network Design," in *3rd International Symposium on Modeling and Optimization in Mobile, Ad Hoc, and Wireless Networks (WiOpt)* April 3-7, 2005. (Acceptance ratio=44/141=31%)
- F.Y.S. Lin and **P.L. Chiu**, "Energy-Efficient Sensor Network Design Subject to Complete Coverage and Discrimination Constraints," in *Second Annual IEEE Communications Society Conference on Sensor and Ad Hoc Communications and Networks (IEEE SECON)*, 26-29 September 2005, pp. 586 - 593. (Acceptance ratio=55/202=27.2%)
- **P.L. Chiu** and F.Y.S. Lin, "Sensor Placement Algorithms Supporting Differentiated Positioning Services," Submitted to *Globecom 2007 Ad-hoc and Sensor Networking Symposium*.

*Dissertation Proposal*

## Sensor Deployment Algorithms for Target Positioning Services

## Q&A