# 論文進度報告

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## Title

- Recovery and Resource Reallocation Strategies to Maximize Network Survivability for Multi-Stage Defense Resource Allocation under Malicious Attacks
- 考量惡意攻擊情況下多階段防禦資源分配以最大化 網路存活度之修復與資源重分配策略

## Agenda

- Problem Description
- Problem Formulation
- Solution Approach

## **Problem Description**

#### **Problem Description**

- Role
  - Defender
  - Attacker
- The network survivability is measured by **average DOS**.

### Defender

• Objective

The defender tried to minimize the damage of the network (Average DOS).

- Budget Constraint (reallocating & new allocated budget)
  - deploying the defense budget in nodes
  - repairing the compromised node

#### Attacker

• Objective

The attacker tried to maximize the damage of the network (Average DOS)

- Budget Constraint
  - deploying the attack budget in nodes

#### Scenario In Each Round

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## Scenario (Defender)

withdraw the resources





#### reallocating & new allocated budget







## **Problem Formulation**

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## **Problem Assumption**

- 1. The problem involves attacker and defender.
- 2. Both attacker and defender have complete information about the network topology.
- 3. Both attacker and defender are limited by budget.
- 4. Only node attack is considered. (Link attack is not considered)
- 5. Only malicious attack is considered. (We do not consider random error)
- 6. The attacker can accumulate experience.

## **Problem Assumption**

- 7. For the defender, the budget can be reallocated and the discount factor is considered.
- 8. For the defender, the compromised node can be repaired.
- 9. Only static network is considered. (We do not consider the growth of network overtime)
- 10. The network survivability is measured by average DOS.
- 11. Any two nodes of the network can form an OD pair.
- 12. We determined the attack success probability using by contest success function, considering the resource allocation of both parties.

## Given

- 1.The network topology
- 2.Attacker's total budget
- 3.Defender's total budget

## Objective

• To minimize the maximum damage of the network (i.e. the average DOS)

## Subject To

- Budget constraint for attacker
- Budget constraint for defender

#### To Determine

- Attacker
  - How to allocate attack budget to each node in each round

- Defender
  - How to allocate defense budget to each node in each round
  - Whether to repair the compromised node in each round

## **Given Parameter**

Given par	Given parameter					
Notation	Description					
V	Index set of nodes					
R	Index set of rounds in the attack and defense actions					
Â	Total budget of attacker					
Â	Total budget of defender					
$ heta_i$	Existing defense resource allocated on node <i>i</i> , where $i \in V$					
e <sub>i</sub>	Repair cost of defender when node <i>i</i> , is dysfunctional, where $i \in V$					
d <sub>ri</sub>	The discount rate of defender reallocate resources on node <i>i</i> , where $i \in V$ and $r \in R$					

### **Decision Variable**

Decision variable					
Notation	Description				
$\vec{a}_r$	Attacker's budget allocation, which is a vector of defense resource $a_{rl_i}$				
$\vec{b}_r$	Defender's budget allocation, which is a vector of attack cost $b_{rl_i} b_{r2}$ to $b_{ri_i}$ in round <i>r</i> , where $i \in V$ and $r \in R$ .				
$a_{ri}$	Attacker's budget allocation on node <i>i</i> in round <i>r</i> , where $i \in V$ and $r \in R$ .				
b <sub>ri</sub>	Defender's budget allocation on node <i>i</i> in round <i>r</i> , where $i \in V$ and $r \in R$ .				
$A_r$	Attacker's total budget in round <i>r</i> , where $r \in R$				
$B_r$	Defender's defense budget in round <i>r</i> , where $r \in R$				

### **Decision Variable**

#### **Decision variable**

Notation	Description
$\overline{z_r}$	Defender's recovery budget allocation, which is a vector of repaired status $z_{r1}$ , $z_{r2}$ to $z_{ri}$ in round $r$ , where $i \in V$ and $r \in R$
Z <sub>ri</sub>	1 if node $i$ is repaired by defender in round $r$ , 0 otherwise where and $i \in V$ and $r \in R$
$\overline{D}(\overline{a}_r, \overline{b}_r)$	The average DOS, which is considering under attacker's and defender's budget allocation are $\vec{a}_r$ and $\vec{b}_r$ in round $\gamma$ , where $r \in R$

#### Formulation

#### **Objective function:**

 $\min_{\overline{\vec{b}_r},\overline{z_r}}\max_{\overline{\vec{a}_r}}\overline{D}(\vec{a}_r,\vec{b}_r)$ 

#### Subject to:

$\sum_{i \in V} b_{ri} + \sum_{i \in V} e_{ri} z_{ri} \le B_r + \sum_{i \in V} \theta_i d_{ri}$	$\forall r \in R$	(IP 1.1)
$\sum_{i \in V} a_{ri} \le A_r$	$\forall r \in R$	(IP 1.2)
$\sum_{r \in R} B_r \le \hat{B}$		(IP 1.3)
$\sum_{r=R} A_r \le \hat{A}$		(IP 1.4)

(IP 1)

## **Solution Approach**

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#### Gradient

- In vector calculus, the gradient of a scalar field is a vector field which points in the direction of the greatest rate of increase of the scalar field, and **whose magnitude is the greatest rate of change**. (from Wikipedia)
- We could use the **partial deviation** to get the increase rate of each variable.
  - ex. F(X,Y) = 2X + 3Y $\frac{\partial F(X,Y)}{\partial X} = 2$   $\frac{\partial F(X,Y)}{\partial Y} = 3$

#### Gradient

• It is almost impossible to calculate the partial-deviation of **Average DOS of each node**, because it is too complicated.

• So....

We calculate the partial-deviation by

$$\lim_{h \to 0} \frac{\overline{D}(r_i + h) - \overline{D}(r_i)}{h}$$

 $r_i$  means the resources on node *i* 

## Example1

• There are three nodes in the network. (AC link is an O-D pair)



	А	В	С
defender resource	1	1	1
attacker resource	1	1	1
Contest intensity(m)	1		

成功機率

0.50

0.50

0.50

#### 平均 DOS =1.5

А	В	С	機率(P)	DOS	P*DOS
0	0	0	0.125	0	0
0	0	1	0.125	1	0.125
0	1	0	0.125	1	0.125
0	1	1	0.125	2	0.25
1	0	0	0.125	1	0.125
1	0	1	0.125	2	0.25
1	1	0	0.125	2	0.25
1	1	1	0.125	3	0.375
				平均DOS	1.5

## Experience

	А	В	С
defender resource	2	1	1
attacker resource	1	1	1
Contest intensity(m)	1		

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А	В	С	機率(P)	DOS	P*DOS
0	0	0	0.166667	0	0
0	0	1	0.166667	1	0.166667
0	1	0	0.166667	1	0.166667
0	1	1	0.166667	2	0.333333
1	0	0	0.083333	1	~~333
1	0	1	0.083333	2	0.166667
1	1	0	0.083333	2	2
1	1	1	0.083333	3	0.25
				平均DOS	1.333333

平均 DOS =1.33333

	А	В	С	
defender resource	1	2	1	
attacker resource	1	1	1	
contest intensity(m)	1			
		平均 DC	OS =1.33333	0.166667
				2010/11

	А	В	С	
defender resource	1	1	2	
attacker resource	1	1	1	
contest intensity(m)	1			
		平均 D	OS = 1.33333	0.166667

### Result1

• Therefore we could know that the important degree of node 1, 2, 3 are same for defender, we need to put the same amount of resources into those node.

### Example2

• There are three nodes in the network. (AC link is an O-D pair)



	А	В	С
defender resource	1	1	1
attacker resource	1	1	1
Contest intensity(m)	1		



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	А	В	С	
defender resource	2	1	1	
attacker resource	1	1	1	
contest intensity(m)	1			
		平均 DO	oS = 0.833333	0.166667

	А	В	С		
defender resource	1	2	1		
attacker resource	1	1	1		
contest intensity(m)	1				
平均 DOS = 1					

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	А	В	С		
defender resource	1	1	2		
attacker resource	1	1	1		
contest intensity(m)	1				
0.166667 平均 DOS = 0.833333					

#### Result2

• Therefore we could know that the node 1,3 are more important than node 2, we could recycle the resource of node 2 and distribute those resources to the node 1,3.



## Gradient

• As a result, we could use the concept of gradient to distribute resources more effectively.

### Thank you for your listening !

## Average DOS Calculating Method

- The more number of the nodes of network, the more difficult to calculate the Average DOS.
- Because we need to consider all of the possible network statuses,

Number of nodes	Network Status
4	$2^4$
10	$2^{10}$
100	$2^{100}$

## Average DOS Calculating Method

- In order to calculate the Average DOS, we need to calculate the **DOS value** and **probability** of each kind of network status.
- However, in order to calculate DOS value of each kind of network status, it need to take lots of time to calculate those value.

#### Example

 There are five nodes of the network and there exists two O-D pairs in the network(AE and BD)



#### • As a result, we need to consider 64 possible network statuses.